

A General n-Port Network's Maximum Transfer-Power Theorem

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Abstract

In this paper, a general n-port network's maximum transfer-power theorem has been proved in detail with a special method by applying a general n-port network's equivalent voltage sources theorem which has been presented with this paper at the same time.

Keywords

Equivalent Voltage-Source, Network, Linear Time-Invariant, Port-Current, Port-Voltage

普遍n口网络的最大传输功率定理

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摘要

本文由作者推导出的普遍n口网络的等效电压源定理详细地推导出普遍n口网络的最大传输功率定理的传输功率的表示式, 并用特殊方法证明传输最大功率的条件。

关键词

等效电压源, 网络, 线性时不变, 端口电流, 端口电压

1. 前言

此问题虽有人研究过[1] [2], 但缺少详细概念论述和严格系统证明。本文根据作者早年在另一篇论文中推导出来的普遍 n 口网络的等效电压源定理, 直接求出端口电流, 再详细推导出传输功率的表示式, 并用特殊方法证明传输最大功率的条件及传输最大功率时的传输效率。不仅研究结果对供电系统有理论价值, 而且对非电系统的模拟研究亦有参考意义。

2. 端口电流

图 1 中 N_s 为由线性时不变元件组成的含有正弦电源的 n 口连通网络, N_L 为由线性时不变元件组成的无源 n 口连通网络, 由 N_s 传输电功率给 N_L , 端口的电流及端口电压的参考方向如图 1 所示。由作者推导出的普遍 n 口网络的等效电压源定理[3], 得出端口电流

$$\dot{i} = (Z_L + Z_S)^{-1} \dot{U}_0 \quad (1)$$

式中 Z_L 及 Z_S 分别表示无源网络 N_L 及 N_0 的端口开路阻抗矩阵, 而 N_0 为网络 N_s 内的电源不作用时的无源网络。 $\dot{U}_0 = [\dot{U}_{01}, \dot{U}_{02}, \dots, \dot{U}_{0n}]^T$ 为 N_s 的端口开路电压(见图 2), $\dot{i} = [\dot{i}_1 \dots \dot{i}_n]^T$ 为端口电流。

关于 \dot{U}_0 , Z_L 及 Z_S 的求得, 在作者的普遍 n 口含源网络的等效电压源定理[3], 此次投稿的另一篇论文)中已有详细的介绍, 故不再赘述, 计算如图 3 所示。

3. 传输功率的表示式

设 1 个端口电压为 $\dot{U}_1 = U_1 e^{j\theta_u}$, 端口电流为 $\dot{i}_1 = I_1 e^{j\theta_i}$, 则 $\dot{U}_1^* = U_1 e^{-j\theta_u}$, $\dot{i}_1^* = I_1 e^{-j\theta_i}$, 角符号“*”意为共轭(conjugate)。

$$\begin{aligned} \dot{U}_1^* \dot{i}_1 &= U_1 e^{-j\theta_u} I_1 e^{j\theta_i} = U_1 I_1 e^{-j(\theta_u - \theta_i)} = U_1 I_1 e^{-j\varphi_1} \\ &= U_1 I_1 \cos \varphi_1 - j U_1 I_1 \sin \varphi_1 = P_1 - j Q_1 \\ \dot{i}_1^* \dot{U}_1 &= I_1 e^{-j\theta_i} U_1 e^{j\theta_u} = U_1 I_1 e^{j(\theta_u - \theta_i)} = U_1 I_1 e^{j\varphi_1} \\ &= U_1 I_1 \cos \varphi_1 + j U_1 I_1 \sin \varphi_1 = P_1 + j Q_1 \end{aligned}$$

两式左右两边相加, 除以 2, 得一个端口传输的功率为

$$P_1 = \frac{1}{2} [\dot{U}_1^* \dot{i}_1 + \dot{i}_1^* \dot{U}_1], \quad (2)$$

因 n 口的传输功率等于各口的传输功率之和, 得

$$\begin{aligned} P &= P_1 + P_2 + \dots + P_n = \frac{1}{2} [(\dot{U}_1^* \dot{i}_1 + \dot{i}_1^* \dot{U}_1) + (\dot{U}_2^* \dot{i}_2 + \dot{i}_2^* \dot{U}_2) + \dots + (\dot{U}_n^* \dot{i}_n + \dot{i}_n^* \dot{U}_n)] \\ &= \frac{1}{2} [(\dot{U}_1^* \dot{i}_1 + \dot{U}_2^* \dot{i}_2 + \dots + \dot{U}_n^* \dot{i}_n) + (\dot{i}_1^* \dot{U}_1 + \dot{i}_2^* \dot{U}_2 + \dots + \dot{i}_n^* \dot{U}_n)] \\ &= \frac{1}{2} [\dot{U}^{*T} \dot{i} + \dot{i}^{*T} \dot{U}] = \frac{1}{2} [\dot{U}^+ \dot{i} + \dot{i}^+ \dot{U}] \end{aligned} \quad (3)$$

式中 $\dot{U}^+ = \dot{U}^{*T} = [\dot{U}_1^* \ \dot{U}_2^* \ \dots \ \dot{U}_n^*]^T$, $\dot{U} = [\dot{U}_1 \ \dot{U}_2 \ \dots \ \dot{U}_n]^T$,

$$\dot{i}^+ = \dot{i}^{*\text{T}} = [\dot{i}_1^* \ \dot{i}_2^* \ \dots \ \dot{i}_n^*]^\text{T}, \quad \dot{i} = [\dot{i}_1 \ \dot{i}_2 \ \dots \ \dot{i}_n]^\text{T},$$

矩阵右上角的符号“+”表示该矩阵的共轭转置矩阵(conjugate transpose), 设 N_L 网络的端口开路阻抗矩阵

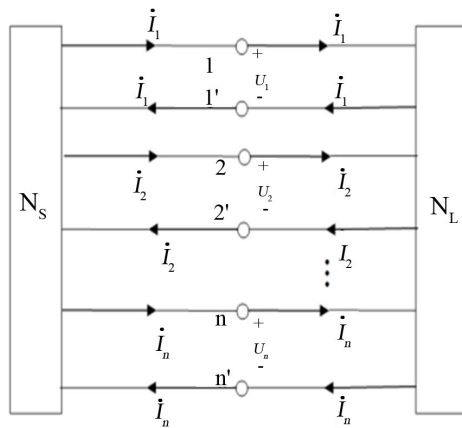


Figure 1. N_S connects N_L
图 1. N_S 连 N_L

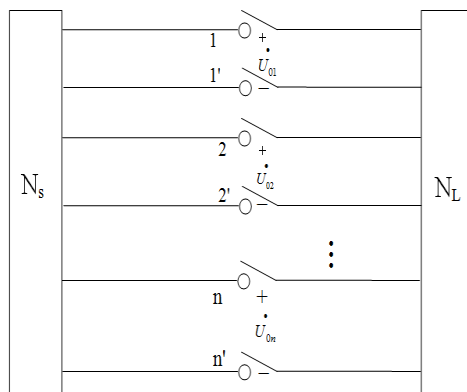


Figure 2. N_S disconnects N_L
图 2. N_S 不连 N_L

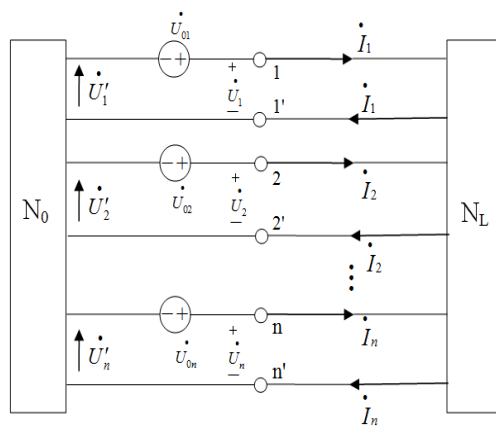


Figure 3. Equivalent voltage sources act on N_0 , N_L
图 3. 等效电压源作用于 N_0 及 N_L

$$Z_L = \begin{bmatrix} Z_{L11} & \cdots & Z_{L1n} \\ \vdots & \ddots & \vdots \\ Z_{Ln1} & \cdots & Z_{Lnn} \end{bmatrix} = \begin{bmatrix} R_{L11} + jX_{L11} & \cdots & R_{L1n} + jX_{L1n} \\ \vdots & \ddots & \vdots \\ R_{Ln1} + jX_{Ln1} & \cdots & R_{Lnn} + jX_{Lnn} \end{bmatrix}$$

则它的共轭端口开路阻抗矩阵

$$Z_L^* = \begin{bmatrix} R_{L11} - jX_{L11} & \cdots & R_{L1n} - jX_{L1n} \\ \vdots & \ddots & \vdots \\ R_{Ln1} - jX_{Ln1} & \cdots & R_{Lnn} - jX_{Lnn} \end{bmatrix}, \quad Z_L + Z_L^* = 2 \begin{bmatrix} R_{L11} & \cdots & R_{L1n} \\ \vdots & \ddots & \vdots \\ R_{Ln1} & \cdots & R_{Lnn} \end{bmatrix} = 2R_L$$

得 N_L 网络的端口的开路电阻矩阵为

$$R_L = \frac{1}{2}(Z_L + Z_L^*) = \begin{bmatrix} R_{L11} & \cdots & R_{L1n} \\ \vdots & \ddots & \vdots \\ R_{Ln1} & \cdots & R_{Lnn} \end{bmatrix}$$

因功率只消耗在电阻上, 对一个电阻 R 而言, 其值为 $P = I^2 R = \dot{I}^* R \dot{I}$, 对矩阵 R_L 而言, 其值为 $P = \dot{I}^+ R_L \dot{I}$, 式中 $\dot{I}^+ = \dot{I}^{*\text{T}} = [\dot{I}_1^* \ \dot{I}_2^* \ \cdots \ \dot{I}_n^*]^{\text{T}}$ 。由此得

$$P = \frac{1}{2} \dot{I}^+ (Z_L + Z_L^*) \dot{I} \quad (4)$$

把(1)式代入(4)式, 得传输给网络 N_L 的功率

$$P = \frac{1}{2} [(Z_L + Z_S)^{-1} \dot{U}_0]^+ (Z_L + Z_L^*) (Z_L + Z_S)^{-1} \dot{U}_0 \quad (5)$$

4. 传输最大功率的条件

因(5)式中, \dot{U}_0 及 Z_S 是不随 Z_L 而变的常量, 故功率 P 只是 Z_L 的函数, 那么, Z_L 在什么条件下, 功率 P 最大呢?

今由(3)式用特殊方法求证。

$$\text{因 } P = \frac{1}{2} [\dot{U}^+ \dot{I} + \dot{I}^+ \dot{U}]$$

$$\Delta P = \frac{1}{2} [\Delta \dot{U}^+ \dot{I} + \dot{U}^+ \Delta \dot{I} + \Delta \dot{I}^+ \dot{U} + \dot{I}^+ \Delta \dot{U}]$$

利用 $\dot{U} = \dot{U}_0 - Z_S \dot{I}$, $\dot{U}^+ = \dot{U}_0^+ - \dot{I}^+ Z_S^+$, \dot{U}_0 及 Z_S 均为常量, $\Delta \dot{U} = -Z_S \Delta \dot{I}$, $\Delta \dot{U}^+ = -\Delta \dot{I}^+ Z_S^+$, 代入, 得

$$\Delta P = \frac{1}{2} [-\Delta \dot{I}^+ Z_S^+ \dot{I} + (\dot{U}_0^+ + \dot{I}^+ Z_S^+) \Delta \dot{I} + \Delta \dot{I}^+ (\dot{U}_0 - Z_S \dot{I}) + \dot{I}^+ (-Z_S \Delta \dot{I})]$$

因 P 最大时, $\Delta P = 0$, $\Delta Z_L \rightarrow 0$, $\Delta \dot{I} \rightarrow 0$, 代入上式, 得

$$0 = -\Delta \dot{I}^+ Z_S^+ \dot{I} + \Delta \dot{I}^+ (\dot{U}_0 - Z_S \dot{I}) = \Delta \dot{I}^+ (-Z_S^+ \dot{I} + \dot{U}_0 - Z_S \dot{I}),$$

$$0 = -Z_S^+ \dot{I} + \dot{U}_0 - Z_S \dot{I},$$

$$\text{即 } (Z_S^+ + Z_S) \dot{I} = \dot{U}_0$$

与(4)式 $(Z_L + Z_S) \dot{I} = \dot{U}_0$ 比较,

$$\text{得 } Z_L = Z_S^+ \quad (6)$$

这就是说当负载阻抗 Z_L 等于电源网络的内阻抗 Z_S 的共轭转置阻抗时, 传输给负载网络的功率最大。这就是最大传输功率定理。

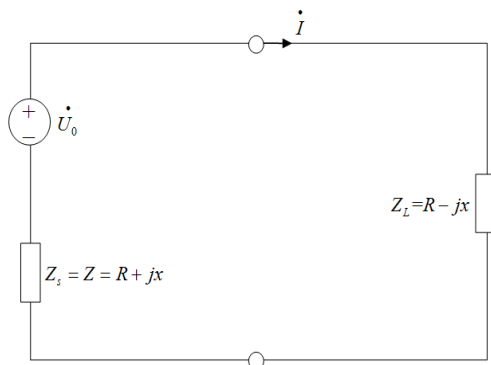


Figure 4. One port network
图 4. 单口网络

5. 传输的最大功率及其传输效率

$$\begin{aligned} \text{由(3)式, } P &= \frac{1}{2} [\dot{U}^+ i + i^+ \dot{U}] = \frac{1}{2} [(Z_L i)^+ i + i^+ (Z_L i)] \\ &= \frac{1}{2} [i^+ Z_L^+ i + i^+ (\dot{U}_0 - Z_s i)] = \frac{1}{2} [i^+ Z_L^+ i + i^+ \dot{U}_0 - i^+ Z_s i] \end{aligned}$$

因功率 P 最大时, $Z_L = Z_s^+$, 即 $Z_L^+ = Z_s$, 故得最大实功率

$$P_m = \frac{1}{2} i^+ \dot{U}_0 \quad (7)$$

网络 N_s 的电源发出的总功率为

$$P_s = \frac{1}{2} [\dot{U}_0^+ i + i^+ \dot{U}_0] = \frac{1}{2} [\dot{U}_0^+ (Z_L + Z_s)^{-1} \dot{U}_0 + i^+ \dot{U}_0] = \frac{1}{2} [i^+ \dot{U}_0 + i^+ \dot{U}_0] = i^+ \dot{U}_0 = 2P_m \quad (8)$$

由(7)式及(8), 可知功率最大时输电效率为 50%, 另一半的功率消耗在电源网络 N_s 的电阻上。

注意: 在上面的证明中, 由 $i = (Z_L + Z_s)^{-1} \dot{U}_0$, $i^+ = \dot{U}_0^+ [(Z_L + Z_s)^{-1}]^+ = \dot{U}_0^+ (Z_L + Z_s)^{-1}$, 为什么 $[(Z_L + Z_s)^{-1}]^+ = (Z_L + Z_s)^{-1}$? 因网络 N_L 及 N_0 纯由电阻(R), 电感(L 及 M)及电容(C)组成, 它们的阻抗矩阵 Z_L 及 Z_s 都是对称的。最大传输功率时, 因 $Z_L = Z_s^+$, 故 $Z_L + Z_s = Z_s^+ + Z_s$ 为纯电阻的对称矩阵, 逆矩阵亦是, 而纯电阻对称矩阵的共轭转置矩阵等于原来的矩阵, 故

$$[(Z_L + Z_s)^{-1}]^+ = (Z_L + Z_s)^{-1}$$

6. 单口网络特例

如图 4, 这时的 Z_L 及 Z_s 各为一个复数, 可看成为 1×1 阶的矩阵, 转置就是它本身, 所以在单口网络情况下, 当 $Z_L = Z_s^+$ 时, 即负载阻抗与电源的内阻抗成共轭复数时, 输出的功率最大。

设 $Z_s = R + jX$, $Z_L = R - jX$, 电源的有效值为 U_0 , 则电流有效值的大小为

$$I = \frac{U_0}{\sqrt{(R+R)^2 + (X-X)^2}} = \frac{U_0}{2R}, \text{ 传输的最大功率为 } P_m = I^2 R = \left(\frac{U_0}{2R}\right)^2 R = \frac{U_0^2}{4R}, \text{ 而电源产生的功率}$$

$$P_0 = U_0 I = U_0 \frac{U_0}{2R} = \frac{U_0^2}{2R} = 2P_m。$$

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